Balanced Trees - 2-3 Tree

Balanced Trees

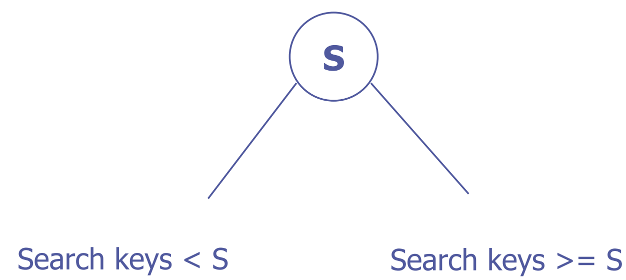
- Binary search trees are not guaranteed to be balanced given random inserts and deletes
  - Tree could degrade to O(n) operations
- Balanced search trees
  - Operations maintain a balanced tree

2-3 Tree

- Guaranteed to always be balanced
  - O(lg n) operations
- Each interior node has two or three children
  - Nodes with 2 children are called 2 nodes
  - Nodes with 3 children are called 3 nodes
  - NOT A BINARY TREE
- Data is stored in both internal nodes and leaves

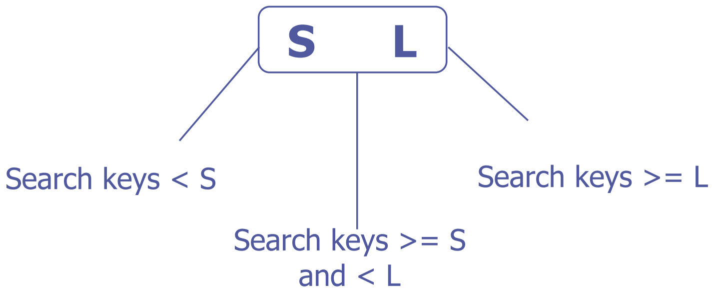
2 Node

- 2 nodes have one data item and 2 children



3 Node

- 3 nodes have two data items and 3 children (a left child, a middle child, and a right child)

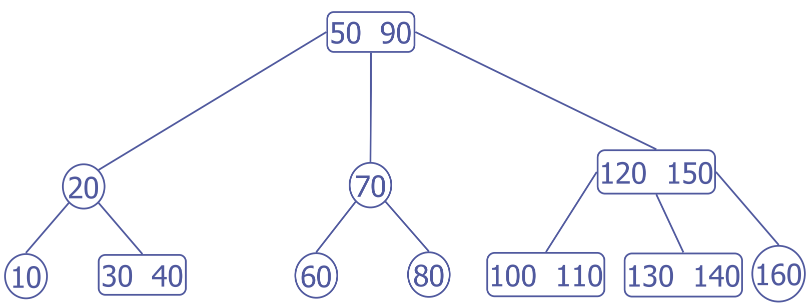


2-3 Tree

- A leaf may contain 1 or 2 data items
- 2-3 trees are good because they are easy to maintain as balanced
  - Operations take care of that for you

Node class
itemtype smallItem, largeItem
*Node *left, *middle, *right, *parent*
// you may find that you need other data members
// to make your life easier
// do not add a third item or a fourth child

2-3 Tree



Traversing a 2-3 Tree

- Inorder traversal -

inorder (node cur)*
if current
inorder(cur->left)
visit small item if it exists
inorder(cur->middle)
visit large item if it exists
inorder(cur->right)

Searching a 2-3 Tree

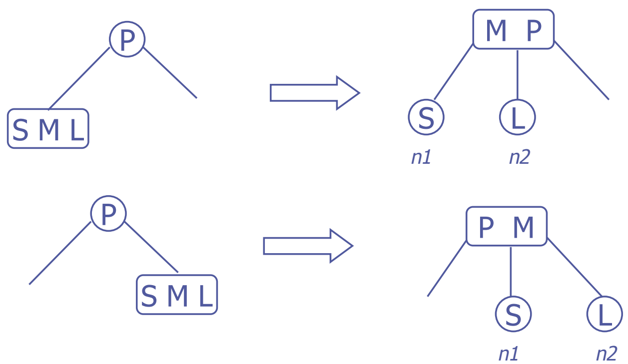
// Assumes small and large exist. You will need to modify
// to account for nodes with only one value

```
search (Node* cur, itemtype key)
    if (cur)
        if (key is in cur)
            return cur
        else
            if (key < cur->small)
                search down left child
            else if (key > cur->large)
                search down right child
            else
                search down middle child
```


Insert

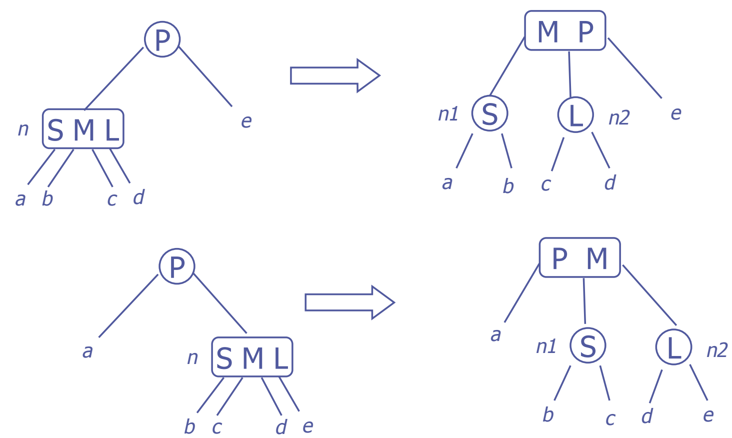
- To insert an item, find a leaf to put the item in then split nodes if necessary

Splitting a Leaf

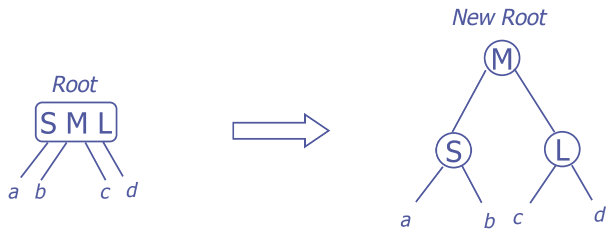


If splitting node causes the parent to have 3 items and 4 children, you will then split an internal node…

Splitting an Internal Node

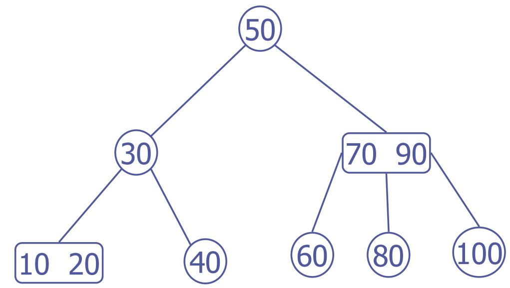


Splitting the Root



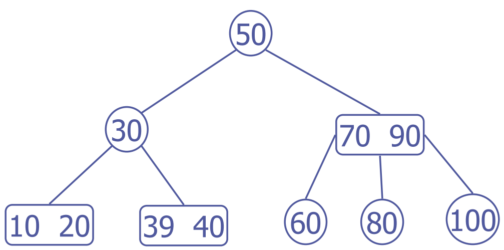
2-3 Tree

Start with this tree



Insert 39

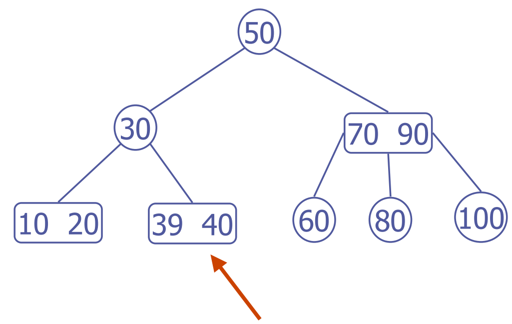
- Locate leaf to insert 39



- Leaf to insert only has 1 data item
  - Add 39 to the leaf

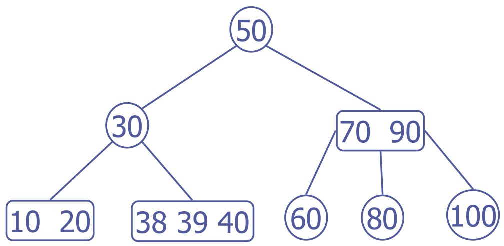
Insert 38

- Locate leaf to insert 38



Insert 38

- Conceptualize inserting 38 into this leaf
  - Do not actually add the item because the node can only hold 2 data items

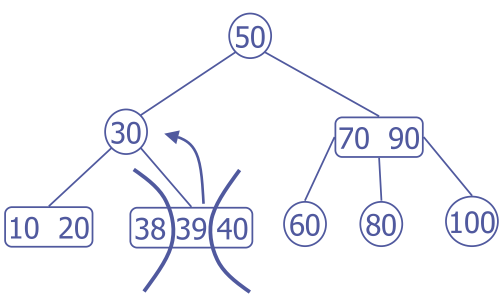


Insert 38

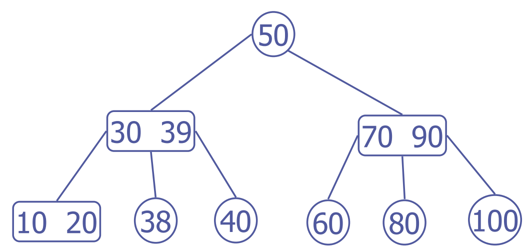
- Determine
  - Smallest = 38
  - Middle = 39
  - Largest = 40

Insert 38

- Move middle value up to parent p
- Separate small and large values into two separate nodes that will be children of p

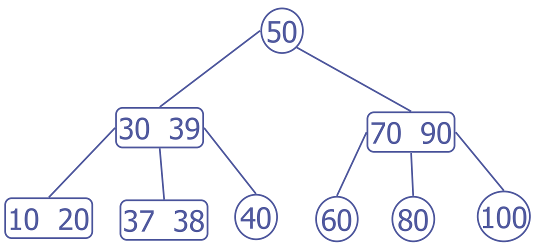


Insert 38



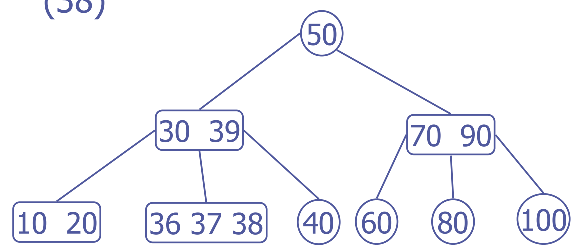
Insert 37

- Locate leaf to insert 37
- Leaf contains 1 data value, just insert value



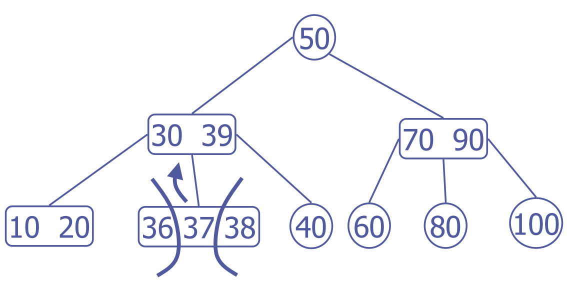
Insert 36

- Locate leaf to insert 36
- Conceptualize inserting 36 into this leaf
  - Determine small (36), middle (37), and large (38)



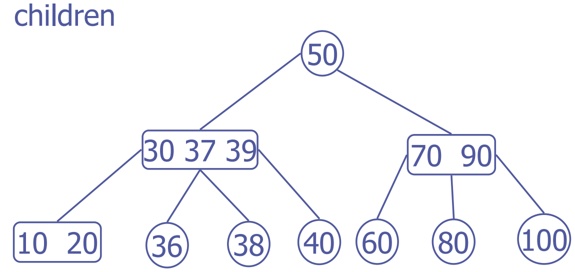
Insert 36

- Conceptualize moving middle value up to parent p
  - Do not actually move, node can't have 3 data values



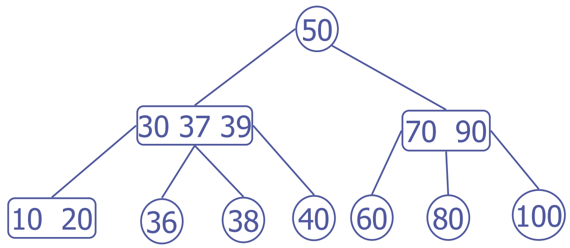
Insert 36

- Conceptualize attaching as children to p the smallest and largest values
  - Do not actually attach because a node can't have 4 children



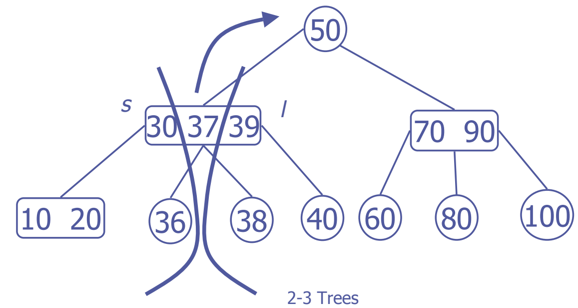
Insert 36

- Parent p now has 3 data values and 4 children
- Split - similar to leaf situation where leaf has 3 data values
  - You can generalize both situations into one



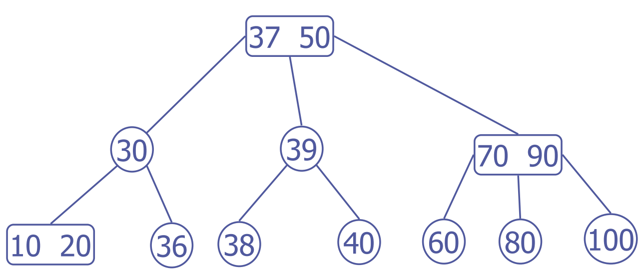
Insert 36

- Split parent p
  - Divide to small (30), middle (37), and large (39)
  - Move middle value to nodes parent
  - Small and large become new children, s and l



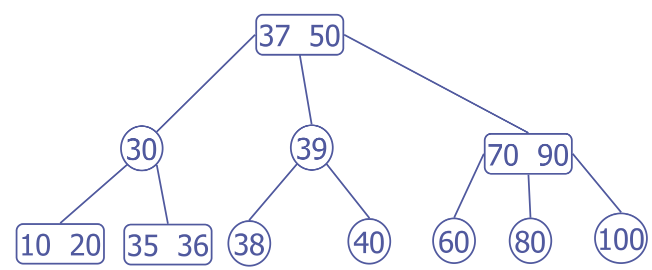
Insert 36

- Divide 4 children
  - Two leftmost become children of s
  - Two rightmost become children of l



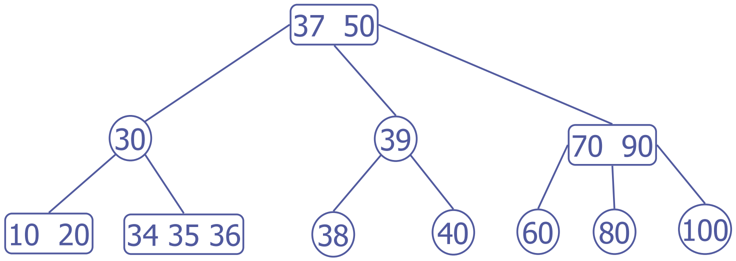
Insert 35

- Insert 35
  - Inserts into leaf



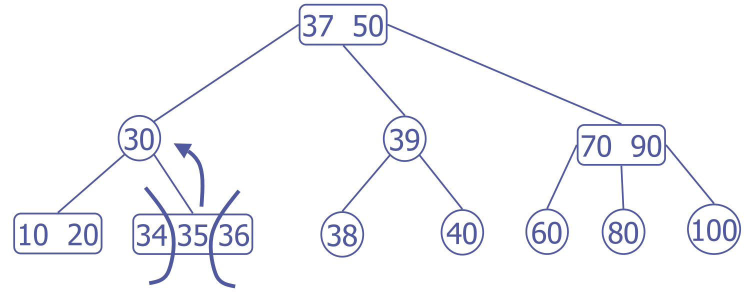
Insert 34

- Insert 34
  - Causes a split



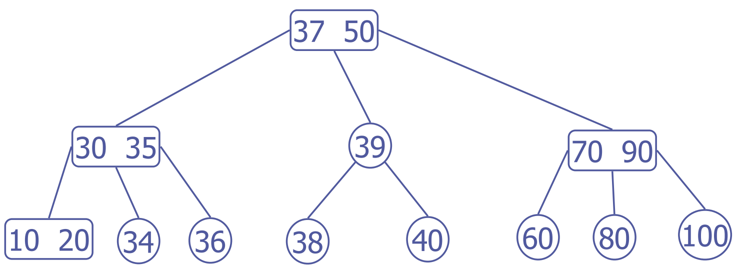
Insert 34

- Insert 34
  - Causes a split



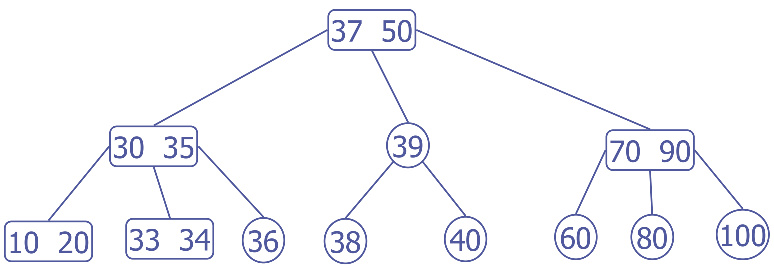
Insert 34

- Insert 34
  - Causes a split



Insert 33

- Insert 33
  - Inserts into leaf



Insert

◆Insert into a tree without duplicates

```
insert (itemtype item)
        leaf = leaf node to insert item (may be null or have 1 or 2 data items)

        if (leaf is null - only happens when root is null)
                add new root to tree with item
        else if (# data items in leaf = 1)
                add item to node
        else // leaf has 2 data items
                split ( leaf, item )
```


Insert (continued)

```
// Item is to be inserted into n.  The insertion
// of item will cause n to have 3 items so n
// must be split
split ( Node* n, itemtype item, … // you may need more )
        if ( n is the root )
                create a new node p
        else
                let p be the parent of n

        Replace node n with 2 nodes, n1 and n2, so that p is their parent

        Give n1 the item in n with the smallest value
        Give n2 the item in n with the largest value

        // continued on next slide…
```

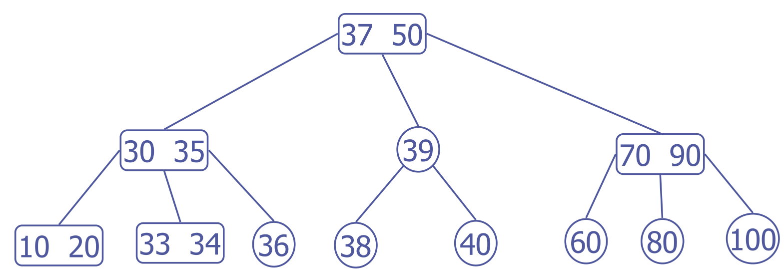

Insert (continued)

```
if (n is not a leaf)
        n1 becomes the parent of n's two leftmost children
        n2 becomes the parent of n's two rightmost children

x = the item in n that has the middle key value

if ( adding x to p would cause p to have 3 items )
        split (p, x)
else
        add x to p
```


Insert 32

◆In class exercise

- Insert 32 into the tree below



Insert 32

Insert 32

Insert 32

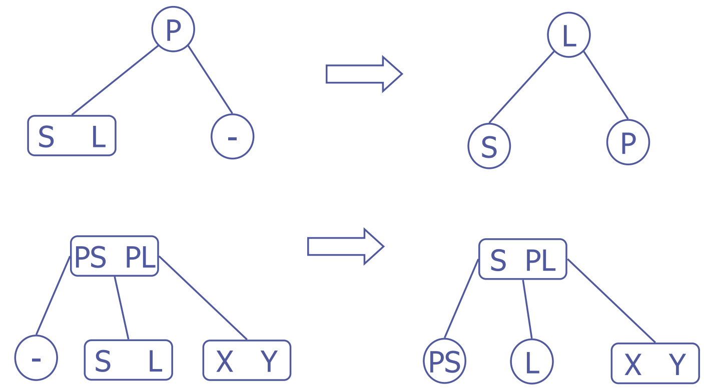
Insert 32

Remove

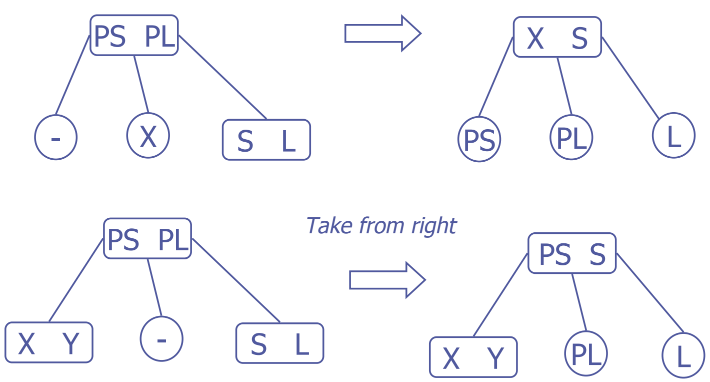
- With insertion, we split nodes. With removing, we merge nodes
- Deletion process needs to begin with a leaf but you might be deleting a value that is not a leaf
  - Swap item to delete with inorder successor

Delete 70

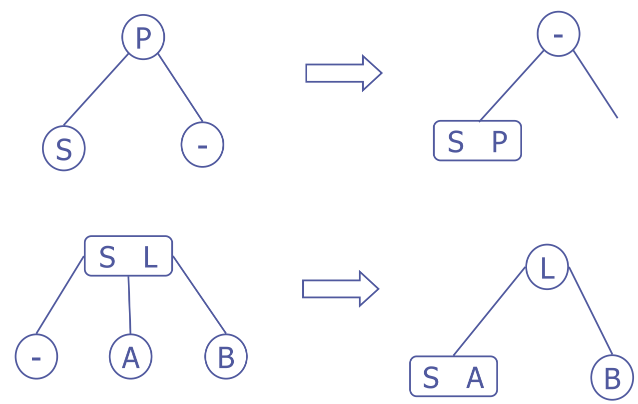
Remove - Redistribute



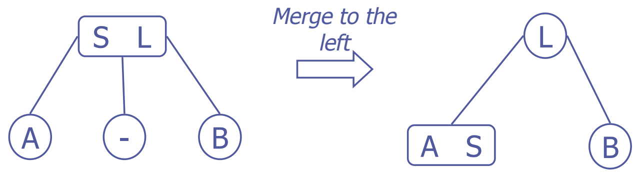
Remove - Redistribute



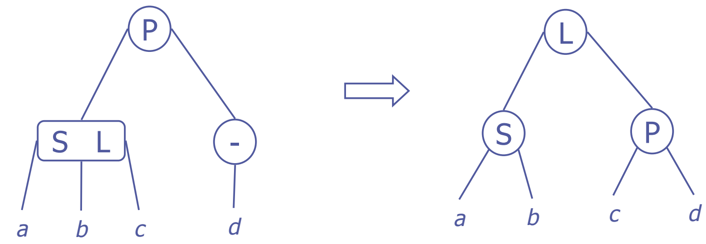
Remove - Merge



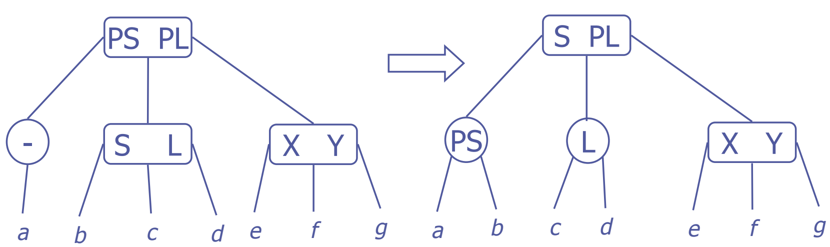
Remove - Merge



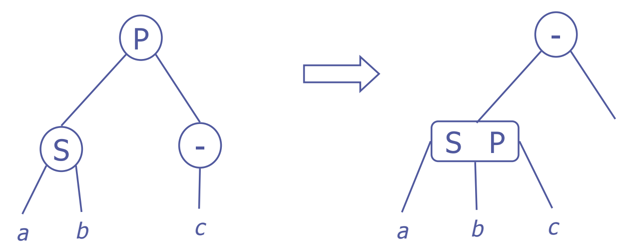
Remove - Redistribute



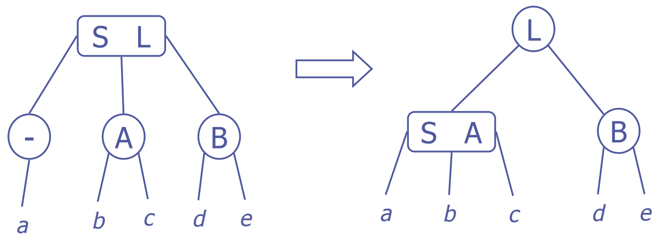
Remove - Redistribute



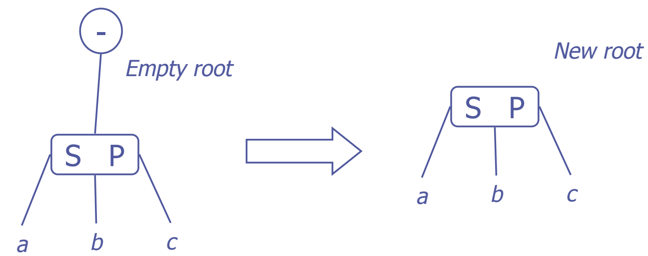
Remove - Merge



Remove - Merge

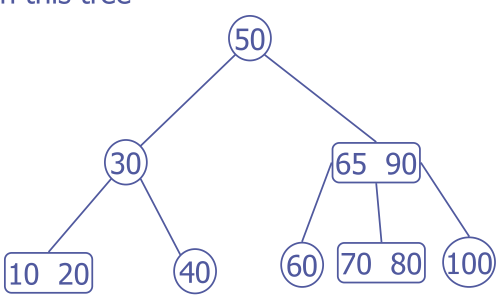


Deleting the Root



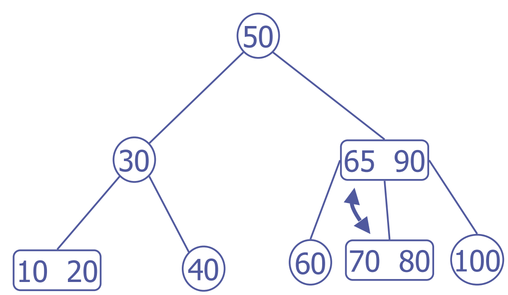
2-3 Tree

Start with this tree



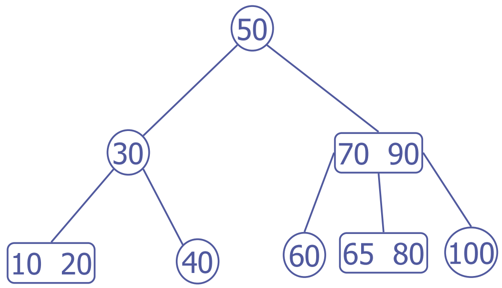
Remove 65

- 65 is an internal node - swap with inorder successor
  - Inorder successor will always be in a leaf



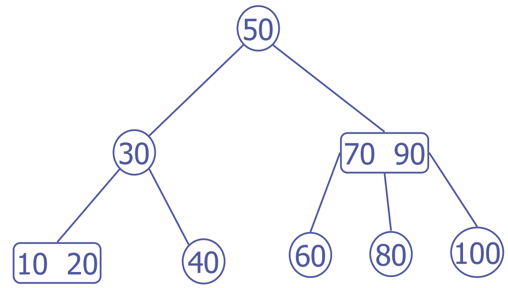
Remove 65

◆65 is now in an invalid location but that is okay because we will remove it



Remove 65

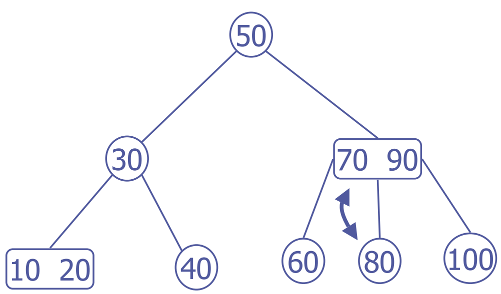
◆Since there are 2 data values in the leaf, just remove data value



Delete 70

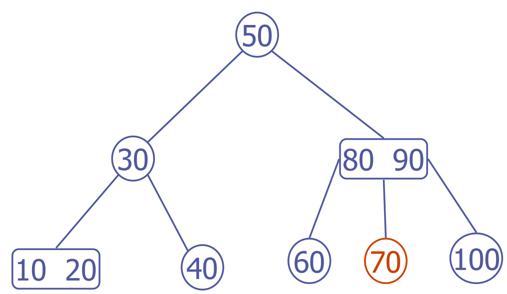
◆70 is an internal node - swap with inorder successor

- Inorder successor will always be in a leaf



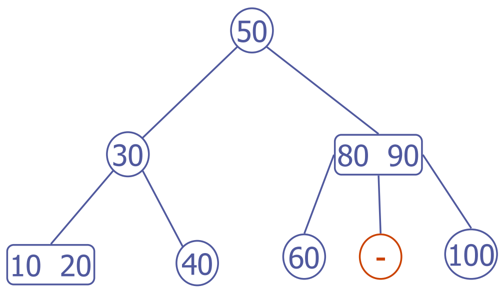
Delete 70

◆70 is now in an invalid location but that is okay - we will be removing that node



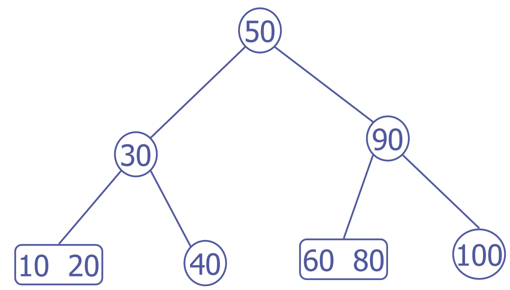
Delete 70

- Removing leaf leaves us with an invalid 2-3 tree



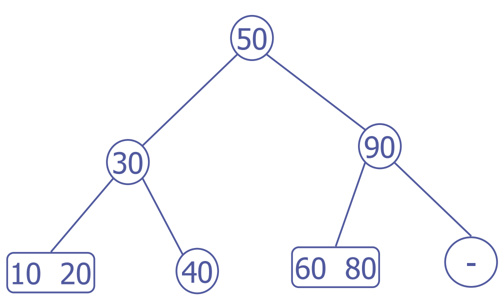
Delete 70

- Merge nodes to fix tree



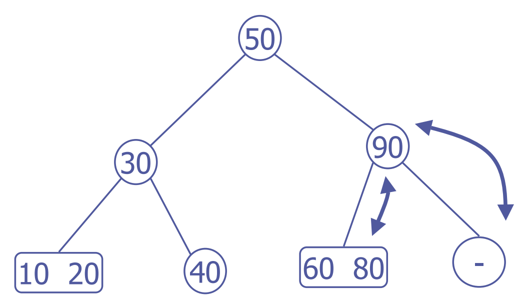
Delete 100

- 100 is already leaf, just remove leaf



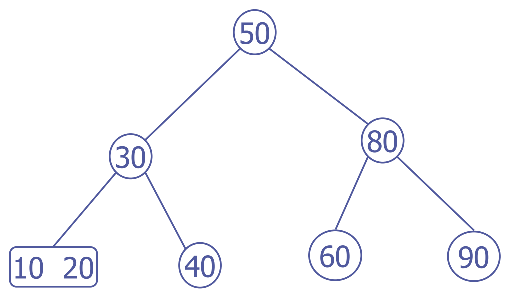
Delete 100

- Sibling has data item to spare, redistribute



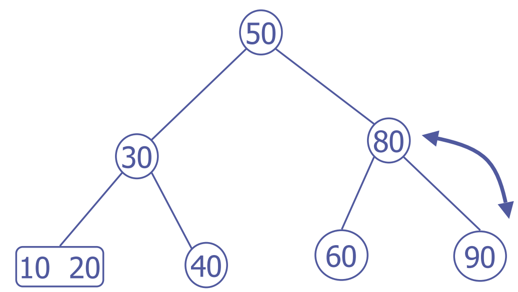
Delete 100

- Sibling has data item to spare, redistribute



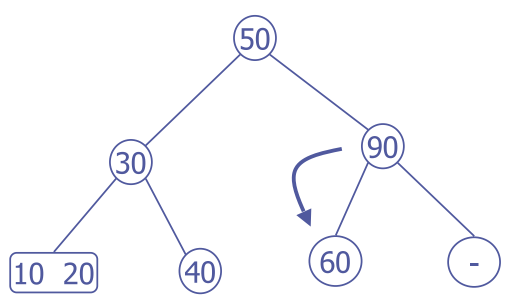
Delete 80

- Swap 80 with inorder successor



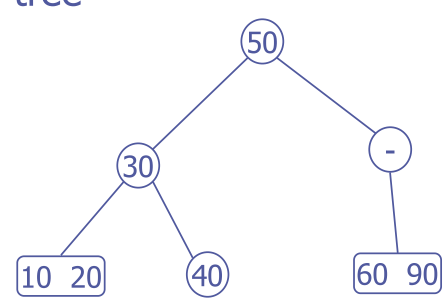
Delete 80

- Can't redistribute so merge nodes



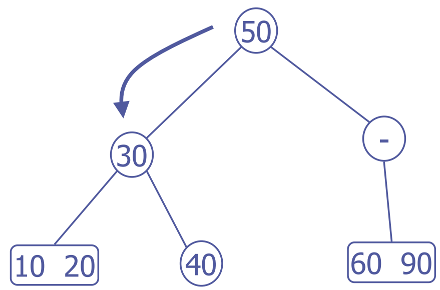
Delete 80

- Can't redistribute so merge nodes
- Invalid 2-3 tree, continue recursively up the tree



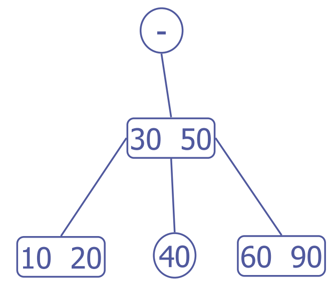
Delete 80

◆Can't redistribute so merge nodes



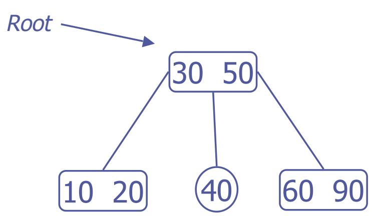
Delete 80

◆Can't redistribute so merge nodes



Delete 80

◆Root is now empty, set new root pointer



Delete

deleteItem (itemtype item)

node = node where item exists (may be null if no item)

if (node)

if (item is not in a leaf)

swap item with inorder successor (always leaf)

leafNode = new location of item to delete

else

leafNode = node

delete item from leafNode

if (leafNode now contains no items)

fix (leafNode)

Delete

// completes the deletion when node n is empty by either
// removing the root, redistributing values, or merging nodes.
// Note: if n is internal, it has only one child

```
fix (Node* n, …) //may need more parameters {
        if (n is the root) {
                remove the root
                set new root pointer
        }
        else {
                Let p be the parent of n
```


Delete

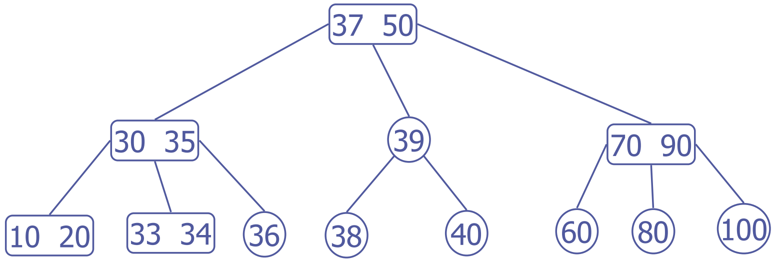
```
if ( some sibling of n has two items ) {
        Distribute items appropriately among n, the
        sibling and the parent (remember take from
        right first)

        if ( n is internal ) {
                Move the appropriate child from sibling n
                (May have to move many children if
                distributing across multiple siblings)
}
```


Delete

```
        else { //merge nodes
                Choose an adjacent sibling s of n (remember,
                merge left first)
                Bring the appropriate item down from p into s

                if ( n is internal )
                        move n's child to s

                Remove node n

                if ( p is now empty )
                        fix (p)
        } //end if
}//end if
```


Delete

- In class exercise - remove the following values from the tree in this order
  - 37, 70



Delete 37